

# COST Action IC1205 on Computational Social Choice: STSM Report

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I visited Prof. Faliszewski at the AGH University of Science and Technology and worked on two projects with him: Shift Bribery under non-monotone rules, and Combinatorial Voter Controls.

The input to the Shift Bribery problem consists of an election  $E = (C, V)$ , where  $C$  is a set of alternatives and  $V = (v_1, v_2, \dots, v_n)$  is a set of voters, a preferred candidate  $p \in C$ , a list  $\Pi = (\pi_1, \pi_2, \dots, \pi_n)$  of price functions, and a budget  $B \in \mathbb{N}$ . Each voter has a preference order which is a linear order over  $C$ . Each price function  $\pi_i : [m] \rightarrow \mathbb{N}$  is a non-decreasing function specifying the cost of shifting  $p$  forward by a given number of positions in the preference order of voter  $v_i \in V$ . The problem asks whether there is a list  $\vec{s} := (s_1, s_2, \dots, s_n)$  of shifts which does make  $p$  the winner when applied to  $E$ , and whose cost  $\Pi(\vec{s})$  does not exceed  $B$ . During my visit, we studied this problem for non-monotone voting rules such as STV, Greedy Monroe, and Greedy Chamberlin-Couran. We found that Shift Bribery for STV is  $W[2]$ -hard for the parameter “total number of shifts”.

The second problem is a new model we developed which occurs naturally in the context of election control problems. In this setting, we are given an election  $E = (C, V)$ , a preferred alternative  $p \in C$ , a list  $W$  of “unregistered” voters, a bundling function  $\kappa : W \rightarrow 2^W$ , and an integer bound  $k \in \mathbb{N}$ , and ask whether there is a size-at-most- $k$  subset  $W' \subseteq W$  such that  $p$  wins the election  $(C, V + \kappa(W'))$ . We studied this problem for two voting rules, Plurality and Condorcet, for different parameters. It turns out that while in general such combinatorial control is computationally significantly harder than its non-combinatorial version, there are some naturally tractable cases. We submit our results to the MFCS 2014.

Finally, I thank the STSM committee for giving me the chance to pay a visit to Prof. Faliszewski. I very much enjoyed staying there and working with him.