# Lattices of Intermediate Theories via Ruitenburg's Theorem\*

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#### 1. Introduction

We propose a systematic study of a class of lattices of intermediate theories stemming from Ruitenburg's Theorem, among which we can find the lattice of negative variants of intermediate logics.

Negative variants ([3, 7]) are theories obtained from intermediate logics by adding the clause  $\neg \neg p \leftrightarrow p$  and closing under modus ponens, thus requiring that the  $\vee$ -free fragment of the logic behaves classically. *Inquisitive logic* InqB ([5]), characterizable as the negative variant of Medvedev's logic, is an example of such logics extensively studied in the literature ([3, 4, 7]).

Building on Ruitenburg's Theorem and a recent algebraic approach to InqB ([1]), we define the lattices of  $\chi$ -variants, generalizing negative variants, and develop a semantic framework to study them. In particular, we will develop an algebraic semantics for these lattices, based on the algebraic semantics introduced in [1]. For each lattice we will also define a suitable concept of variety, obtaining a correspondence result in the style of Birkhoff.

As an application of this framework, we will show that negative logics extending InqB are linearly ordered under inclusion. To prove this, we will define and study a suitable generalization of locally-finite varieties in the framework of  $\chi$ -variants.

## 2. $\chi$ -variants

Given an intermediate logic L, its negative variant is defined as the set  $L^{\neg} := \{\phi(\overline{p}) \mid \phi[\neg \overline{p}/\overline{p}] \in L\}^1$ . As shown in [3]§3.2, negative variants can be axiomatized by a special class of Hilbert-style derivation systems, since  $L^{\neg} = \text{MP}(L \cup \{\neg \neg p \leftrightarrow p \mid p \text{ atomic proposition}\})^2$ —compare with Definition 3.2.13 of [3]. This identity justifies the intuition that atoms behave like  $\neg \neg$ -fixpoints for negative variants.

A surprising result by Ruitenburg allows us to generalize the concept of negative logic, following the fixpoint intuition. Let  $\chi(p)$  be a formula in one propositional variable. Ruitenburg's Theorem ([8]; see also [6, 9] for a recent semantic proof) states that the sequence  $p, \chi(p), \chi^2(p) := \chi(\chi(p)), \ldots$  is definitely periodic with period 2, modulo logical equivalence; that is, there exists an n such that  $\vdash \chi^n(p) \leftrightarrow \chi^{n+2}(p)$ , where  $\vdash$  is intuitionistic provability. We will call the smallest n for which this holds the *index of*  $\chi$ .

We define the  $\chi$ -variant of an intermediate logic L as  $L^{\chi} := \{\phi(\overline{p}) \mid \phi[\chi^n(\overline{p})/\overline{p}] \in L\}$ , for n the index of  $\chi$ . Intuitively, we are requiring atomic proposition to play the role of  $\chi^2$ -fixpoints, in analogy with the case of negative variants. And in fact, we obtain an analogous axiomatization result:  $L^{\chi} = MP(L \cup \{\chi^2(p) \leftrightarrow p \mid p \text{ atomic proposition}\})$ .

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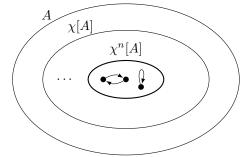
<sup>&</sup>lt;sup>1</sup>With  $\overline{p}$  we denote a sequence of variables  $p_1, \ldots, p_k$ . Given  $\chi(q)$  a formula with only one free variable, we denote by  $\chi(\overline{p})$  the sequence  $\chi(p_1), \ldots, \chi(p_k)$ . We denote by  $\phi[\overline{\psi}/\overline{p}]$  the formula obtained by substituting every instance of  $p_i$  with the formula  $\psi_i$ .

 $<sup>^2\</sup>mathrm{MP}$  denotes the closure under modus ponens.

# 3. Algebraic semantics for $\chi$ -variants

An algebraic semantics for InqB was presented in [1]. Following the same approach, we can define a semantics for  $\chi$ -variants. The key to this generalization lies in an algebraic interpretation of Ruitenburg's Theorem.

Given an Heyting algebra A, we can treat  $\chi$  as an endomorphism of A, mapping  $a \in A$  to  $[\![\chi(p)]\!]_{[p \mapsto a]}$ . Ruitenburg's Theorem tells us exactly that  $\chi$  restricted to the set  $A^{\chi} = \chi^n[A]$  is an *involution*, and so that  $A^{\chi}$  are exactly the  $\chi^2$ -fixpoints of the algebra. Notice that, in the case of negation,  $A^{\neg p}$  is exactly the set of regular elements of the algebra.



We call a valuation<sup>3</sup>  $\sigma: AP \to A$  a  $\chi$ -valuation if the range of  $\sigma$  is contained in  $A^{\chi}$ . We can define a semantics suitable to study  $\chi$ -variants simply by restricting the valuations we consider to be  $\chi$ -valuations: we say that  $\phi$  is  $\chi$ -valid on A if  $\llbracket \phi \rrbracket_{\sigma}^{A} = \top$  for every  $\chi$ -valuation  $\sigma$ . We will denote by  $\mathrm{Th}^{\chi}(A)$  the set of valid formulas in A; for  $\mathcal C$  a class of Heyting algebras, we will denote by  $\mathrm{Th}^{\chi}(A) = \bigcap \{\mathrm{Th}^{\chi}(A) \mid A \in \mathcal C\}$ .

A routine check shows that  $\operatorname{Th}^{\chi}(\mathcal{C}) = (\operatorname{Log}(\mathcal{C}))^{\chi}$ . In particular, every class of algebras defines a  $\chi$ -variant via the semantics introduced; and every negative logic  $L^{\chi}$  is the class of some class of Heyting algebras, since  $\operatorname{Th}(\operatorname{Var}(L)) = L^4$ . Thus for every  $\chi$ -variant  $L^{\chi}$ , the semantics introduced is *correct and complete* if restricted to the appropriate class of Heyting algebras.

This result can be further refined, obtaining a correspondence between  $\chi$ -variants and some special varieties, in the same style as Birkhoff Theorem (Theorem 11.9 in [2]). The key observation is that, under a  $\chi$ -valuation, the semantics of a formula in A is restricted to the sub-algebra generated by  $A^{\chi}$ . In case this sub-algebra coincides with the whole algebra, we say that A is core-generated.

If we denote by  $V_{CG}$  the set of core-generated elements of the variety V, the observation above entails that  $V_{CG}$  determines univocally  $\operatorname{Th}^{\chi}(V)$ . Moreover, if  $V_{CG} = V'_{CG}$  then  $\operatorname{Th}^{\chi}(V) = \operatorname{Th}^{\chi}(V')$ . This suggests the definition of the following closure operator on varieties:  $\operatorname{Cl}^{\chi}(V) := \{B \mid \exists A \in V. A \leq B \text{ and } A^{\chi} = B^{\chi}\}^{5}$ —notice that  $V_{CG} = V'_{CG}$  iff  $\operatorname{Cl}^{\chi}(V) = \operatorname{Cl}^{\chi}(V')$ . Indeed, it is easy to verify that given a  $\chi$ -variant  $L^{\chi}$ , the set  $\operatorname{Var}^{\chi}(L^{\chi}) := \{A \mid L^{\chi} \subseteq \operatorname{Th}^{\chi}(A)\}$  is a variety closed under the operation  $\operatorname{Cl}^{\chi}$ .

**Theorem 1.** Let  $\mathbf{IL}^{\chi}$  and  $\mathbf{HA}^{\chi}$  denote the lattice of  $\chi$ -variants and the lattice of varieties closed under  $\mathrm{Cl}^{\chi}$  respectively. The function  $\mathsf{Var}^{\chi}: \mathbf{IL}^{\chi} \to \mathbf{HA}^{\chi}$  is an isomorphism and  $\mathrm{Th}^{\chi} = (\mathsf{Var}^{\chi})^{-1}$ .

$$\begin{array}{c}
\mathbf{IL} \xrightarrow{\bullet^{\chi}} \mathbf{IL}^{\chi} \\
\mathsf{Var} \downarrow & & \downarrow \mathsf{Var}^{\chi} \\
\mathbf{HA} \xrightarrow{\operatorname{Cl}^{\chi}} \mathbf{HA}^{\chi}
\end{array}$$

<sup>&</sup>lt;sup>3</sup>AP denotes the set of atomic propositions.

 $<sup>^4</sup>$ With Var(L) we denote the variety of Heyting algebras validating all formulas in L.

<sup>&</sup>lt;sup>5</sup>With  $A \leq B$  we denote that A is a subalgebra of B.

### 4. Application: extensions of IngB

As an application of the framework introduced, we prove that negative logics extending InqB are linearly ordered by inclusion. By Theorem 1, this implies that subvarieties of  $\mathsf{Var}^{\neg p}(\mathsf{InqB})$  are linearly ordered. An important feature, that we can exploit to prove this, is that coregenerated elements of  $\mathsf{Var}^{\neg p}(\mathsf{InqB})$  are locally finite—compare with Lemmas 4.1 and 4.3 of [1]. This property corresponds to the usual locally-finiteness for varieties of Heyting algebras, and can be studied more in general for varieties in  $\mathsf{HA}^{\chi}$ .

We say a variety  $V \in \mathbf{H}\mathbf{A}^{\chi}$  is  $\chi$ -locally-finite if all the algebras in  $V_{CG}$  are locally-finite; or equivalently, given  $A \in V$  and finitely many elements  $a_1, \ldots, a_k \in A^{\chi}$ , the sub-algebra  $\langle a_1, \ldots, a_k \rangle_A$  is finite. Notice that, under a  $\chi$ -valuation, the interpretation of a formula  $\phi$  lies in the subalgebra generated by the core, so this seems a sensible generalization of locally-finiteness to study  $\chi$ -logics and the corresponding varieties. And in fact we have the following:

**Lemma 2.** If V is  $\chi$ -locally-finite, then V is generated by its finite, subdirectly-irreducible, core-generated elements.

So, since  $\mathsf{Var}^{\neg p}(\mathsf{InqB})$  is  $\neg p$ -locally-finite, to determine its subvarieties it suffices to consider its finite, subdirectly-irreducible, core-generated elements. These are exactly (up to isomorphism) the finite inquisitive algebras H(B) studied in [1], for B a finite Boolean algebra. In particular, denoting with  $B_k$  the unique (up to isomorphism) Boolean algebra with k generators, we have  $H(B_k) \in \mathsf{Var}(\{H(B_l)\})$  for  $k \leq l$ .

Combining these observations with Theorem 1, we obtain that the negative variants extending properly InqB are all of the form InqB<sub>N</sub> := Th<sup>¬p</sup>({ $H(B_k) \mid k \leq N$ }) for  $N \in \mathbb{N}$ . Moreover, these are all distinct theories since  $C_N \in \text{InqB}_N \setminus \text{InqB}_{N+1}$  for the following:

$$C_N := \bigvee_{\substack{i \neq j \\ 1 \leq i, j \leq N+1}} (p_i \to p_j)$$

Therefore, we obtain the following characterization:

**Theorem 3.** The sublattice of negative variants extending InqB forms the chain:

$$InqB \subsetneq ... \subsetneq InqB_3 \subsetneq InqB_2 \subsetneq InqB_1 = \mathbf{CPC}.$$

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