# **Computational Social Choice: Spring 2017**

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## Plan for Today

The broad aim for today is to show how we can *characterise* voting rules in terms of their properties. We review three approaches:

- Axiomatic method: to characterise a (family of) voting rule(s) as the only one satisfying certain axioms
- *Maximum likelihood estimation:* to characterise a voting rule as computing the most likely "correct" winner, given *n* distorted copies of an objectively "correct" ranking (the ballots)
- *Distance-based rationalisation:* to characterise a voting rule in terms of a notion of *consensus* (profiles where outcomes are clear) and a notion of *distance* (from such a consensus profile)

Under the first approach we think of voting a a *compromise-seeking* activity (so we need to be fair, etc.). Under the second approach we think of voting as a *truth-finding* activity (e.g., amongst experts).

#### **Reminder: Formal Framework**

Need to choose from a finite set  $X = \{x_1, \ldots, x_m\}$  of *alternatives*.

Let  $\mathcal{L}(X)$  denote the set of all strict linear orders on X. We use elements of  $\mathcal{L}(X)$  to model (true) *preferences* and (declared) *ballots*. Each member of a finite set  $N = \{1, \ldots, n\}$  of *voters* supplies us with a ballot, giving rise to a *profile*  $\mathbf{R} = (R_1, \ldots, R_n) \in \mathcal{L}(X)^n$ .

A voting rule (or social choice function) for N and X selects one or more winners for every such profile:

$$F: \mathcal{L}(X)^n \to 2^X \setminus \{\emptyset\}$$

If  $|F(\mathbf{R})| = 1$  for all profiles  $\mathbf{R} \in \mathcal{L}(X)^n$ , then F is called *resolute*. If F is resolute, we usually write  $F(\mathbf{R}) = x^*$  instead of  $F(\mathbf{R}) = \{x^*\}$ . <u>Notation</u>: Write  $N_{x \succ y}^{\mathbf{R}} = \{i \in N \mid (x, y) \in R_i\}$  for the set of voters who rank alternative x above alternative y in profile  $\mathbf{R}$ .

#### **Axioms: Anonymity and Neutrality**

Two basic fairness requirements for a voting rule F:

- *F* is *anonymous* if  $F(R_1, \ldots, R_n) = F(R_{\pi(1)}, \ldots, R_{\pi(n)})$  for any profile  $(R_1, \ldots, R_n)$  and any permutation  $\pi : N \to N$ .
- *F* is *neutral* if  $F(\pi(\mathbf{R})) = \pi(F(\mathbf{R}))$  for any profile  $\mathbf{R}$  and any permutation  $\pi : X \to X$  (with  $\pi$  extended to profiles and sets of alternatives in the natural manner).
- <u>Thus:</u> A is symmetry w.r.t. voters. N is symmetry w.r.t. alternatives.

## **Axiom: Positive Responsiveness**

A (not necessarily resolute) voting rule satisfies *positive responsiveness* if, whenever some voter raises a (possibly tied) winner  $x^*$  in her ballot, then  $x^*$  will become the *unique* winner. Formally:

▶ *F* is *positively responsive* if  $x^* \in F(\mathbf{R})$  implies  $\{x^*\} = F(\mathbf{R'})$ for any alternative  $x^*$  and any two *distinct* profiles  $\mathbf{R}$  and  $\mathbf{R'}$ with  $N_{x^* \succ y}^{\mathbf{R}} \subseteq N_{x^* \succ y}^{\mathbf{R'}}$  and  $N_{y \succ z}^{\mathbf{R}} = N_{y \succ z}^{\mathbf{R'}}$  for all  $y, z \in X \setminus \{x^*\}$ .

Thus, this is a monotonicity requirement (we'll see others later on).

## May's Theorem

When there are only *two alternatives*, then all the voting rules we have seen coincide with the *simple majority rule*. Good news:

**Theorem 1 (May, 1952)** A voting rule for two alternatives satisfies anonymity, neutrality, and positive responsiveness if and only if that rule is the simple majority rule.

This provides a good justification for using this rule (arguing in favour of "majority" directly is harder than arguing for anonymity etc.).

K.O. May. A Set of Independent Necessary and Sufficient Conditions for Simple Majority Decisions. *Econometrica*, 20(4):680–684, 1952.

## **Proof Sketch**

Clearly, the simple majority rule satisfies all three properties.  $\checkmark$ Now for the other direction:

Assume the number of voters is *odd* (other case: similar)  $\sim$  no ties.

There are two possible ballots:  $x \succ y$  and  $y \succ x$ .

Anonymity  $\sim$  only *number of ballots* of each type matters.

Consider all possible profiles R. Distinguish two cases:

• Whenever  $|N_{x\succ y}^{\mathbf{R}}| = |N_{y\succ x}^{\mathbf{R}}| + 1$ , then only x wins.

By *PR*, *x* wins whenever  $|N_{x \succ y}^{\mathbf{R}}| > |N_{y \succ x}^{\mathbf{R}}|$ . By *neutrality*, *y* wins otherwise. But this is just what the simple majority rule does.  $\checkmark$ 

• There exist a profile  $\mathbf{R}$  with  $|N_{x\succ y}^{\mathbf{R}}| = |N_{y\succ x}^{\mathbf{R}}| + 1$ , yet y wins. Suppose one x-voter switches to y, yielding  $\mathbf{R'}$ . By PR, now only y wins. But now  $|N_{y\succ x}^{\mathbf{R'}}| = |N_{x\succ y}^{\mathbf{R'}}| + 1$ , which is symmetric to the earlier situation, so by *neutrality* x should win. Contradiction.  $\checkmark$ 

### Young's Theorem

Another seminal result (which we won't discuss in detail here) is known as *Young's Theorem*. It provides a characterisation of the *PSR*'s. The core axiom is *reinforcement* (a.k.a. *consistency*):

► F satisfies reinforcement if, whenever we split the electorate into two groups and some alternative were to win for both groups, then it will also win for the full electorate. More precisely:

 $F(\mathbf{R}) \cap F(\mathbf{R'}) \neq \emptyset \implies F(\mathbf{R} \oplus \mathbf{R'}) = F(\mathbf{R}) \cap F(\mathbf{R'})$ 

Young showed that F is a (generalised) positional scoring rule iff it satisfies anonymity, neutrality, reinforcement, and a technical condition known as continuity.

Here, "generalised" means that the scoring vector need not be decreasing.

H.P. Young. Social Choice Scoring Functions. *SIAM Journal on Applied Mathematics*, 28(4):824–838, 1975.

## Voting as Truth-Tracking

An alternative interpretation of "voting":

- There exists an objectively "correct" ranking of the alternatives.
- The voters want to identify the correct ranking (or winner), but cannot tell with certainty which ranking is correct. Their ballots reflect what they believe to be true.
- We want to estimate the most likely ranking (or winner), given the ballots we observe. Can we use a voting rule to do this?

## Example

Consider the following scenario:

- two alternatives: x and y
- either  $x \succ y$  or  $y \succ x$  (we don't know which and have no priors)
- 20 voters/experts with probability 75% each of getting it right

Now suppose we observe that 12/20 voters say  $x \succ y$ . What can we infer, given this observation (let's call it E)?

• Probability for E to happen in case  $x \succ y$  is correct:

$$P(E \mid x \succ y) = \binom{20}{12} \cdot 0.75^{12} \cdot 0.25^{8}$$

• Probability for E to happen in case  $y \succ x$  is correct:

$$P(E \mid y \succ x) = \binom{20}{8} \cdot 0.75^8 \cdot 0.25^{12}$$

<u>Thus:</u>  $P(E \mid x \succ y)/P(E \mid y \succ x) = 0.75^4/0.25^4 = 81.$ From Bayes and assuming uniform priors  $[P(x \succ y) = P(y \succ x)]$ : Given E, x being better is 81 times as likely as y being better.

#### The Condorcet Jury Theorem

For the case of two alternatives, the simple majority rule is the best choice also under the truth-tracking perspective:

**Theorem 2 (Condorcet, 1785)** Suppose a jury of n voters need to select the better of two alternatives and each voter independently makes the correct decision with the same probability  $p > \frac{1}{2}$ . Then the probability that the simple majority rule returns the correct decision increases monotonically in n and approaches 1 as n goes to infinity.

<u>Proof sketch</u>: By the law of large numbers, the number of voters making the correct choice approaches  $p \cdot n > \frac{1}{2} \cdot n$ .

For modern expositions see Nitzan (2010) and Young (1995).

Writings of the Marquis de Condorcet. In I. McLean and A. Urken (eds.), *Classics of Social Choice*, University of Michigan Press, 1995.

S. Nitzan. Collective Preference and Choice. Cambridge University Press, 2010.

H.P. Young. Optimal Voting Rules. J. Economic Perspectives, 9(1):51-64, 1995.

### **Characterising Voting Rules via Noise Models**

For n alternatives, Young (1995) shows that, if the probability of a voter to rank a given pair correctly is  $p > \frac{1}{2}$ , then the voting rule selecting the most likely winner coincides with the *Kemeny rule*.

Conitzer and Sandholm (2005) ask a general question:

► For a given voting rule *F*, can we design a "*noise model*" such that *F* is a *maximum likelihood estimator* for the winner?

H.P. Young. Optimal Voting Rules. J. Economic Perspectives, 9(1):51-64, 1995.

V. Conitzer and T. Sandholm. Common Voting Rules as Maximum Likelihood Estimators. Proc. UAI-2005.

E. Elkind and A. Slinko. Rationalizations of Voting Rules. In F. Brandt et al. (eds.), *Handbook of Computational Social Choice*. CUP, 2016.

#### The Borda Rule as a Maximum Likelihood Estimator

It *is* possible for the Borda rule:

**Proposition 3 (Conitzer and Sandholm, 2005)** If each voter independently ranks the true winner at position k with probability  $\frac{2^{m-k}}{2^m-1}$ , then the maximum likelihood estimator is the Borda rule.

<u>Proof</u>: Let  $r_i(x)$  be the position at which voter *i* ranks alternative *x*.

Probability to observe the actual ballot profile if x is the true winner:

$$\frac{\prod_{i \in N} 2^{m-r_i(x)}}{(2^m - 1)^n} = \frac{2^{\sum_{i \in N} m - r_i(x)}}{(2^m - 1)^n} = \frac{2^{\text{BordaScore}(x)}}{(2^m - 1)^n}$$

Hence, x has maximal likelihood of being the true winner  $\underline{\mathrm{iff}}\ x$  has a maximal Borda score.  $\checkmark$ 

V. Conitzer and T. Sandholm. Common Voting Rules as Maximum Likelihood Estimators. Proc. UAI-2005.

#### **Characterisation via Consensus and Distance**

<u>Recall:</u> Rules such as The *Dodgson* and *Young* compute the "closest" profile with a Condorcet winner and then elect that Condorcet winner. This suggests a general method for defining a voting rule:

- Fix a class of *consensus profiles*: profiles in which there is a clear (set of) winner(s). (And specify *who* wins.)
- Fix a metric to measure the *distance* between two profiles.
- This induces a *voting rule:* for a given profile, find the closest consensus profile(s) and elect the corresponding winner(s).

T. Meskanen and H. Nurmi. Closeness Counts in Social Choice. In M. Braham and F. Steffen (eds.), *Power, Freedom, and Voting*, Springer-Verlag, 2008.

E. Elkind and A. Slinko. Rationalizations of Voting Rules. In F. Brandt et al. (eds.), *Handbook of Computational Social Choice*. CUP, 2016.

# **Notions of Consensus**

Four natural definitions for what constitutes a consensus profile:

- Condorcet Winner: there exists a Condorcet winner  $x (\rightsquigarrow x \text{ wins})$
- *Majority Winner:* there exists an alternative x that is ranked first by an absolute majority of the voters ( $\rightsquigarrow x$  wins)
- Unanimous Winner: there exists an alternative x that is ranked first by all voters (→ x wins)
- Unanimous Ranking: all voters report exactly the same ranking (→ the top alternative in that unanimous ranking wins)

(Other definitions are possible.)

### Ways of Measuring Distance

Two natural definitions of distance between profiles R and R':

Swap distance: minimal number of pairs of adjacent alternatives that need to get swapped to get from R to R'.
Equivalently: distance between two ballots = number of pairs of alternatives with distinct relative ranking (so-called Kendall tau distance); sum over voters to get distance between two profiles.

$$\frac{1}{2} \cdot \sum_{i \in N} \#\{(x, y) \in X^2 \mid \mathbb{1}_{i \in N_{x \succ y}^{R}} \neq \mathbb{1}_{i \in N_{x \succ y}^{R'}}\}$$

• *Discrete distance:* distance between two ballots is 0 if they are the same and 1 otherwise; sum over voters to get profile distance.

$$\#\{i \in N \mid R_i \neq R'_i\}$$

(Other definitions are possible.)

# Examples

Two voting rules for which the standard definition is already formulated in terms of consensus and distance:

- Dodgson Rule = Condorcet Winner + Swap Distance
- Kemeny Rule = Unanimous Ranking + Swap Distance

How about other rules? Borda? Plurality?

Writings of C.L. Dodgson. In I. McLean and A. Urken (eds.), *Classics of Social Choice*, University of Michigan Press, 1995.

J. Kemeny. Mathematics without Numbers. Daedalus, 88:571-591, 1959.

#### **Characterisation of the Borda Rule**

<u>Recall</u>: the Borda rule is the PSR with vector  $(m-1, m-2, \ldots, 0)$ .

**Proposition 4 (Farkas and Nitzan, 1979)** Borda is characterised by the unanimous winner consensus criterion and the swap distance.

<u>Proof sketch</u>: The swap distance between a given ballot that ranks x at position k and the closest ballot that ranks x at the top is k-1. Thus, if voter i ranks x at position k, she gives -(k-1) points to x. This corresponds to the PSR with vector  $(0, -1, -2, \ldots, -(m-1))$ , which is equivalent to the Borda rule.  $\checkmark$ 

<u>Remark:</u> So Dodgson, Kemeny, and Borda are all *rationalisable* via the same notion of distance!

D. Farkas and S. Nitzan. The Borda Rule and Pareto Stability: A Comment. *Econometrica*, 47(5):1305–1306, 1979.

## **Characterisation of the Plurality Rule**

<u>Recall</u>: the plurality rule is the PSR with scoring vector (1, 0, ..., 0). **Proposition 5 (Nitzan, 1981)** *Plurality is characterised by the* 

unanimous winner consensus criterion and the discrete distance.

<u>Proof:</u> Immediate. √

<u>Remark 1:</u> to be precise, Nitzan used a slightly different distance

<u>Remark 2</u>: also works with Majority Winner + discrete distance, but doesn't work with Condorcet Winner or Unanimous Ranking

S. Nitzan. Some Measures of Closeness to Unanimity and their Implications. *Theory and Decision*, 13(2):129–138, 1981.

## Summary

We have seen three approaches to characterising a voting rule:

- as the only rule satisfying certain *axioms*;
- as returning *the most likely "true" winner*, given the noisy signals the voters have received about the "true" ranking; and
- as computing the closest *consensus* profile (w.r.t. some *distance*) with a clear winner and returning that winner.

All three approaches are (potentially) useful

- to better understand particular voting rules;
- to explain why there are so many "natural" voting rules; and
- to help prove general results about families of voting rules.

What next? More applications of the axiomatic method.